

Cuckoo Search-Ant Colony Optimization Based Scheduling in Grid Computing

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- 2 Problem Definition/Research Question
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- We use ant colony optimization for selecting the appropriate and optimal resources.

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Block 1

the transaction processing model considering load balanced scheduling is formulated as follows:

$$\begin{aligned} & \text{Minimize} && L \\ & \text{subject to} && \text{proc}_j \leq \text{makespan}(T), \forall j = 1, \dots, n. \end{aligned} \tag{1}$$

where the constraint states that the processing time (proc_j) of all the transactions on node N_j should lie within its makespan.

Table : Divide the solution to K categories based on random cluster numbers in the solution S_1 .

c_1	c_2	c_3
3	5	1
6	7	2
—	8	4

Table : Calculated cost value for all solutions.

Solution	N_1	N_2	N_3	N_4	N_5	N_6	N_7	N_8	Cost
S_1	3	3	1	3	2	1	2	2	6.5
S_2	3	3	3	1	2	2	3	2	7.0
S_3	3	1	2	1	1	3	2	2	7.5

Table : Calculated fitness value for dataset objects by S_1 , S_2 and S_3 .

Nodes	load	Fitness of S_1	Fitness of S_2	Fitness of S_3
1	3	0.5	0.5	0
2	8	3.0	3.0	3.0
3	5	0	1.5	0
4	2	0	0	0
5	4	0	0	1.0
6	6	0.5	1.0	2.0
7	7	1.5	0	1.0
8	6	1.0	1.0	0.5

Table : Parameters of LBTS_ACO

Parameters	Value
α	0.5
β	0.5
ρ	0.2
d	0.25

LBTS_ACO

- 1: Select the node from the appropriate cluster using Cuckoo Search Algorithm
- 2: $k = 1$ ▷ Iteration starts
- 3: $P_k = 0$ ($P_k \leq 1$) $L_j \leq \frac{n}{2} L_{bs} \leftarrow L_j$
- 4: Update the pheromone_load of selected node ▷ pheromone_load update
- 5: Calculate the quality of the node ▷ Finds quality of the node
- 6: $L^P \leftarrow L_{bs}$
- 7: $k = k + 1$
- 8: $P_k = \frac{\log(k)}{\log(K)}$
- 9: select random nodes N_{r1} and N_{r2}
- 10: find L_{r1} and L_{r2}
- 11: Calculate $\phi(L_j, L_{r1}, L_{r2})$ ($0 \leq \phi(L_j, L_{r1}, L_{r2}) \leq 1.5$) $L_{r1} \leq L_{r2}$
- 12: $L_{bs} \leftarrow L_{r1}$
- 13: $L_j \leftarrow L_{r1}$
- 14: $L_{bs} \leftarrow L_{r2}$
- 15: $L_j \leftarrow L_{r2}$
- 16: Update the pheromone_load of the node ▷ pheromone_load update ↻

LBTS_CSACO algorithm

- 1: Initialize the transaction T_i at node N_j
- 2: Find the load L_j of the node N_j
- 3: make N_j the best-so-far node N_{bs} All the transactions are not scheduled
- 4: Call LBTS_ACO
- 5: Make N_j the best-so-far node
- 6: Increment the load of N_j by 1

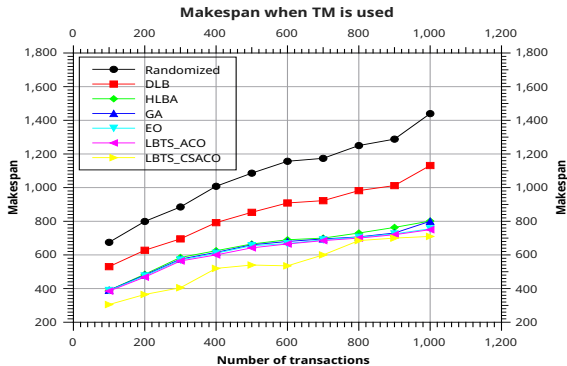


Figure : Makespan with Transaction Management

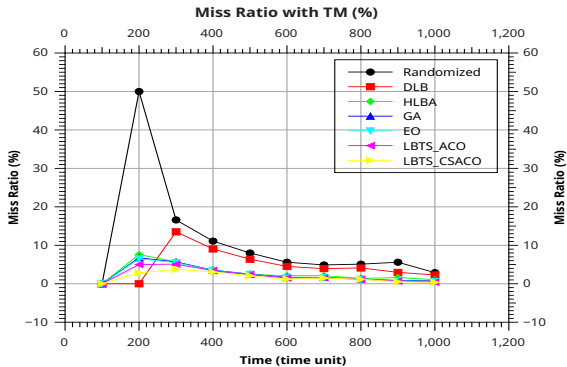


Figure : Miss Ratio with Transaction Management

Thank You